



***10.2 Graph Terminology and  
Special Types of Graphs***

## Objectives:

The main purpose for this lesson is to introduce the following:

- We define Basic Terminology.
- Define the neighborhood and degree of vertex.
- THE HANDSHAKING THEOREM.
- In-degree & Out-degree (directed graph).
- Some Special Simple Graphs.
- Bipartite Graphs.

# *Basic Terminology*

## DEFINITION 1

- Two vertices  $u$  and  $v$  in an undirected graph  $G$  are called adjacent (or neighbors) in  $G$  if  $u$  and  $v$  are endpoints of an edge of  $G$ .
- If  $e$  is associated with  $\{u, v\}$ , the edge  $e$  is called incident with the vertices  $u$  and  $v$ .
- The edge  $e$  is also said to connect  $u$  and  $v$ .
- The vertices  $u$  and  $v$  are called endpoints of an edge associated with  $\{u, v\}$ .

## DEFINITION 2

- The set of all neighbors of a vertex  $v$  of  $G = (V, E)$ , denoted by  $N(v)$ , is called the *neighborhood* of  $v$ . If  $A$  is a subset of  $V$ , we denote by  $N(A)$  the set of all vertices in  $G$  that are adjacent to at least one vertex in  $A$ . So,  $N(A) = \sum_{v \in A} N(v)$ .

## *The degree of a vertex*

### DEFINITION 3

- The degree of a vertex in an undirected graph is the number of edges incident with it, except that a loop at a vertex contributes twice to the degree of that vertex.
- The degree of the vertex  $v$  is denoted by  $\text{deg}(v)$ .
- A vertex of degree zero is called isolated.
- A vertex is pendant if and only if it has **degree one**.

## Example 1:

What are the degrees of the vertices in the graphs  $G$  and  $H$  displayed in Figure 1?

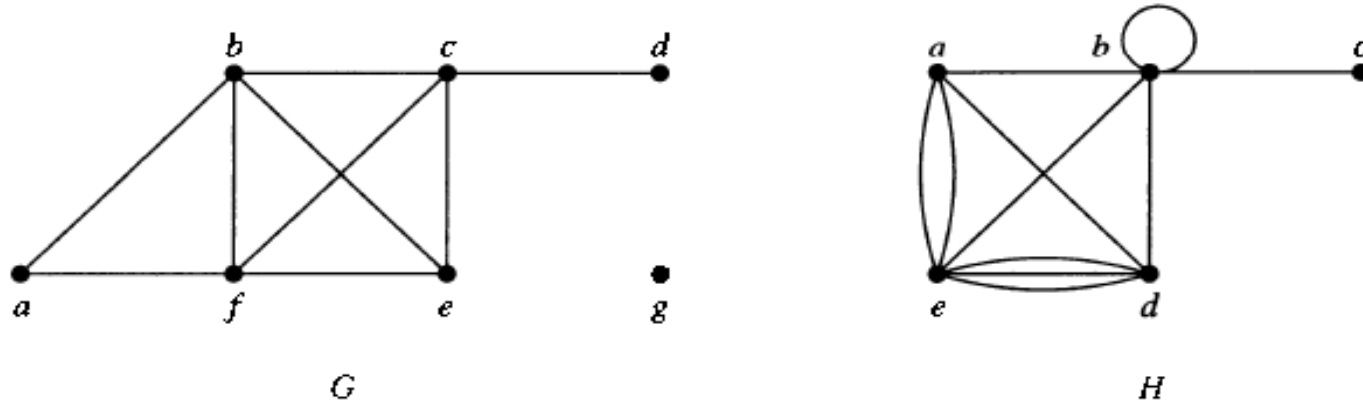


FIGURE 1 The Undirected Graphs  $G$  and  $H$ .

## Solution:

In  $G$ :

$$\deg(a) = 2, \deg(b) = \deg(c) = \deg(f) = 4, \deg(d) = 1, \deg(e) = 3, \text{ and } \deg(g) = 0.$$

- The neighborhoods of these vertices are

$N(a) = \{b, f\}$ ,  $N(b) = \{a, c, e, f\}$ ,  $N(c) = \{b, d, e, f\}$ ,  
 $N(d) = \{c\}$ ,  $N(e) = \{b, c, f\}$ ,  $N(f) = \{a, b, c, e\}$ , and  
 $N(g) = \emptyset$ .

### In H:

$\deg(a) = 4$ ,  $\deg(b) = \deg(e) = 6$ ,  $\deg(c) = 1$ ,  
and  $\deg(d) = 5$ .

The neighborhoods of these vertices are  $N(a) = \{b, d, e\}$ ,  $N(b) = \{a, b, c, d, e\}$ ,  $N(c) = \{b\}$ ,  $N(d) = \{a, b, e\}$ ,  
and  $N(e) = \{a, b, d\}$ .

- A vertex of **degree zero** is called **isolated**.
- It follows that an isolated vertex is not adjacent to any vertex.
- Vertex  $g$  in graph  $G$  in Example 1 is isolated.
- A vertex is **pendant** if and only if it has **degree one**.
- Consequently, a pendant vertex is adjacent to exactly one other vertex.
- Vertex  $d$  in graph  $G$  in Example 1 is pendant.



# ***THE HANDSHAKING THEOREM***

## **THEOREM 1**

Let  $G = (V, E)$  be an undirected graph with  $e$  edges. Then  $2e = \sum_{v \in V} \deg(v)$ .

(Note that this applies even if multiple edges and loops are present.)

### **EXAMPLE 3**

**How many edges are there in a graph with 10 vertices each of degree 6?**

#### **Solution:**

Because the sum of the degrees of the vertices is  $6 \cdot$

$$10 \times 6 = 60,$$

it follows that  $2e = 60$ .

Therefore,  $e = 30$ .

## THEOREM 2

An undirected graph has an even number of vertices of odd degree.

## DEFINITION 4

- When  $(u, v)$  is an edge of the graph  $G$  with directed edges,  $u$  is said to be adjacent to  $v$  and  $v$  is said to be adjacent from  $u$ .
- The vertex  $u$  is called the initial vertex of  $(u, v)$ , and  $v$  is called the terminal or end vertex of  $(u, v)$ .
- The initial vertex and terminal vertex of a loop are the same.

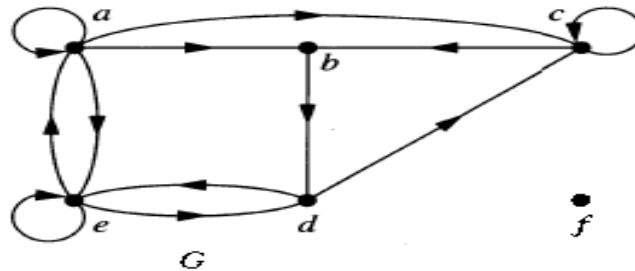
# In-degree & Out-degree (directed graph)

## DEFINITION 4

- In a graph with directed edges the in-degree of a vertex  $v$ , denoted by  $\text{deg}^-(v)$ , is the number of edges with  $v$  as their terminal vertex.
- The out-degree of  $v$ , denoted by  $\text{deg}^+(v)$ , is the number of edges with  $v$  as their initial vertex.
- (Note that a loop at a vertex contributes 1 to both the in-degree and the out-degree of this vertex.)

## EXAMPLE 4

Find the in-degree and out-degree of each vertex in the graph  $G$  with directed edges shown in Figure 2 .



**FIGURE 2** The Directed Graph  $G$ .

### Solution:

The in-degrees in  $G$  are  $\deg^-(a)=2$ ,  $\deg^-(b)=2$ ,  $\deg^-(c)=3$ ,  
 $\deg^-(d)=2$ ,  $\deg^-(e)=3$ , and  $\deg^-(f)=0$ .

The out-degrees are  $\deg^+(a)=4$ ,  $\deg^+(b)=1$ ,  $\deg^+(c)=2$ ,  $\deg^+(d)=2$ ,  
 $\deg^+(e)=3$ , and  $\deg^+(f) = 0$ .

## THEOREM 3

Let  $G = (V, E)$  be a graph with directed edges.

Then:

$$\sum_{v \in V} \deg^-(v) = \sum_{v \in V} \deg^+(v) = |E|.$$

# Some Special Simple Graphs

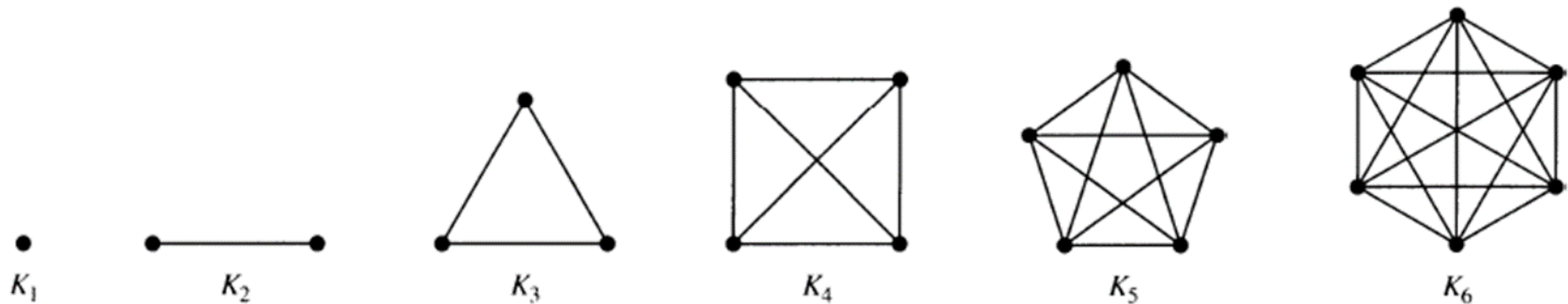
## EXAMPLE 5

*Complete Graphs* The complete graph on  $n$  vertices, denoted by  $K_n$ , is the simple graph that contains exactly one edge between each pair of distinct vertices.

The graphs  $K_n$ , for  $n=1,2,3,4,5,6$ , are displayed in Figure 3.



A simple graph for which there is at least one pair of distinct vertex not connected by an edge is called **noncomplete**



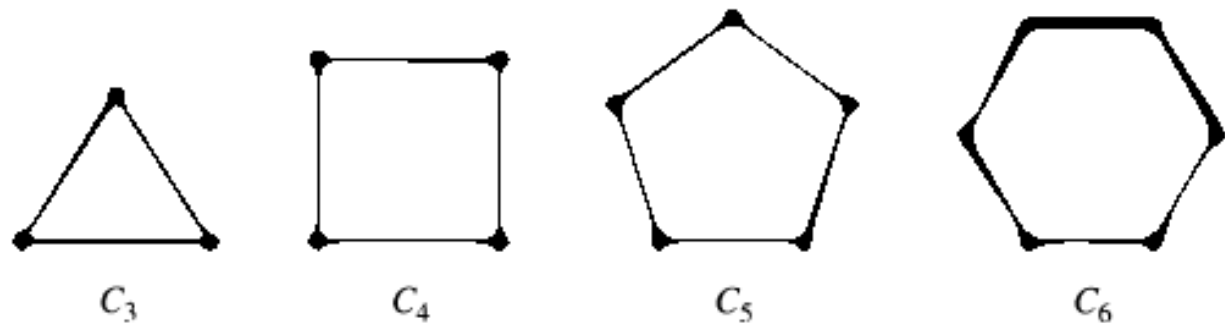
**FIGURE 3** The Graphs  $K_n$  for  $1 \leq n \leq 6$ .

## EXAMPLE 6

Cycles The cycle  $C_n$ ,  $n \geq 3$ , consists of  $n$  vertices  $V_1, V_2, \dots, V_n$  and edges  $\{V_1, V_2\}, \{V_2, V_3\}, \dots, \{V_{n-1}, V_n\}$ , and  $\{V_n, V_1\}$ .

The cycles  $C_3, C_4, C_5$ , and  $C_6$  are displayed in

Figure 4.

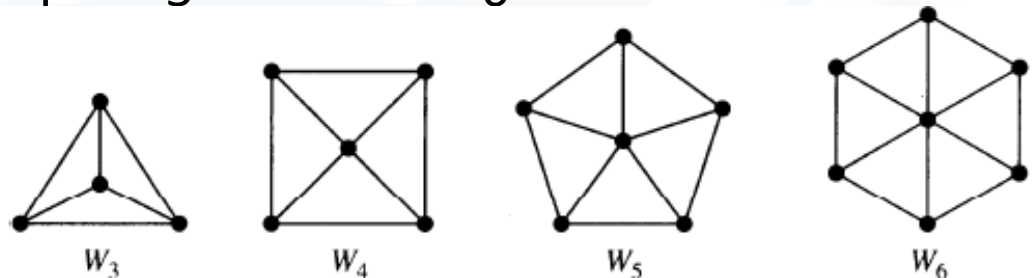


**FIGURE 4** The Cycles  $C_3, C_4, C_5,$  and  $C_6$ .

## EXAMPLE 7

*Wheels* We obtain the wheel  $W_n$  when we add an additional vertex to the cycle  $C_n$ , for  $n \geq 3$ , and connect this new vertex to each of the  $n$  vertices in  $C_n$ , by new edges.

The wheels  $W_3$ ,  $W_4$ ,  $W_5$ , and  $W_6$  are displayed in Figure 5.



**FIGURE 5** The Wheels  $W_3$ ,  $W_4$ ,  $W_5$ , and  $W_6$ .

# How many vertices & edges in each type?

Type of the Simple Graph	Number of Vertices	Number of Edges
$K_n$	$n$	$\frac{n(n-1)}{2}$
$C_n$	$n$	$n$
$W_n$	$n+1$	$2n$

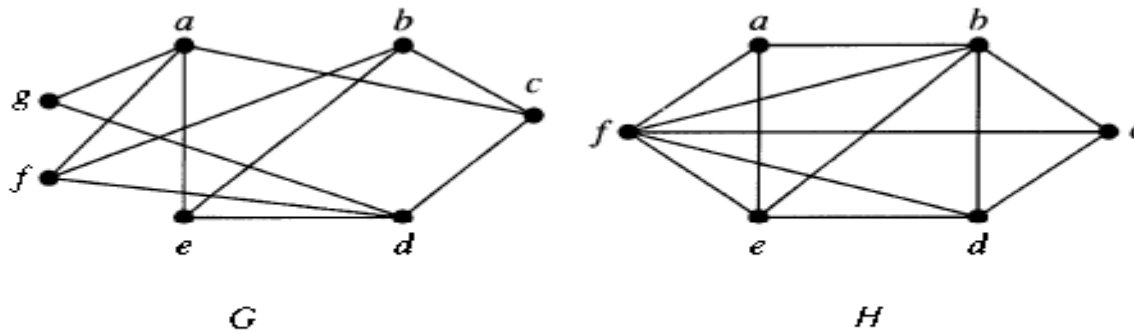
# *Bipartite Graphs*

## DEFINITION 6

A simple graph  $G$  is called bipartite if its vertex set  $V$  can be partitioned into two disjoint sets  $V_1$  and  $V_2$  such that every edge in the graph connects a vertex in  $V_1$  and a vertex in  $V_2$ . (so that no edge in  $G$  connects either two vertices in  $V_1$  or two vertices in  $V_2$ ). When this condition holds, we call the pair  $(V_1, V_2)$  a bipartition of the vertex set  $V$  of  $G$ .

## EXAMPLE 11

Are the graphs  $G$  and  $H$  displayed in Figure 8 bipartite?



**FIGURE 8** The Undirected Graphs  $G$  and  $H$ .

## Solution:

- Graph G is bipartite because its vertex set is the union of two disjoint sets,  $\{a, b, d\}$  and  $\{c, e, f, g\}$ , and each edge connects a vertex in one of these subsets to a vertex in the other subset.

(Note that for G to be bipartite it is not necessary that every vertex in  $\{a, b, d\}$  be adjacent to every vertex in  $\{c, e, f, g\}$ . For instance, b and g are not adjacent.)

- Graph H is not bipartite because its vertex set cannot be partitioned into two subsets so that edges do not connect two vertices from the same subset. verify this by (consider the vertices a, b, and f.)

# *Homework*

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- 1
- 2
- 3
- 7
- 8
- 9
- 10
- 29(a,b,c)
- 37(a,b,d,e,f).