## Logic and Computer Design Fundamentals

# Chapter 5 - Sequential Circuits 

## Part 2 - Sequential Circuit Analysis

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## Overview

- Part 1 - Storage Elements
- Part 2 - Sequential Circuit Analysis
- State tables
- State diagrams
- Equivalent states
- Part 3 - Sequential Circuit Design
- Part 4 - State Machine Design


## Sequential Circuit Analysis

- General Model
- Current State at time $(t)$ is stored in an array of flip-flops.
- Next State at time ( $\mathbf{t}+\mathbf{1}$ ) is a Boolean function of CLK State and Inputs.
- Outputs at time (t) are a Boolean function of State (t) and (sometimes) Inputs (t).


## Example 1 (from Fig. 5-15)

- Input:
$\mathbf{x}(\mathbf{t})$
- Output: $\mathbf{y}(\mathbf{t})$
- State:
(A(t), B(t))
- What is the Output Function?
- What is the Next State Function?



## Example 1 (from Fig. 5-15) (continued)

- Boolean equations for the functions:
- $\mathbf{A}(\mathbf{t}+\mathbf{1})=\mathbf{A}(\mathbf{t}) \mathbf{x}(\mathbf{t})$ $+\mathbf{B}(\mathbf{t}) \mathbf{x}(\mathbf{t})$
- $\mathbf{B}(\mathbf{t} \mathbf{t} \mathbf{1})=\overline{\mathbf{A}}(\mathbf{t}) \mathbf{x}(\mathbf{t})$
- $\mathbf{y}(\mathbf{t})=\overline{\mathbf{x}}(\mathbf{t})(\mathbf{B}(\mathbf{t})+\mathbf{A}(\mathbf{t}))$



## Example 1(from Fig. 5-15) (continued)

- Where in time are inputs, outputs and states defined?
- $\mathbf{A}(\mathbf{t}+\mathbf{1})=\mathbf{A}(\mathbf{t}) \mathbf{x}(\mathbf{t})+\mathbf{B}(\mathbf{t}) \mathbf{x}(\mathbf{t})$
- $\mathbf{B}(\mathbf{t}+\mathbf{1})=\overline{\mathbf{A}}(\mathbf{t}) \mathbf{x}(\mathbf{t})$
- $\mathbf{y}(\mathbf{t})=\overline{\mathbf{x}}(\mathbf{t})(\mathbf{B}(\mathbf{t})+\mathbf{A}(\mathbf{t}))$



## Example 1(from Fig. 5-15) (continued)

- Where in time are inputs, outputs and states defined?



## State Table Characteristics

- State table - a multiple variable table with the following four sections:
- Present State - the values of the state variables for each allowed state.
- Input - the input combinations allowed.
- Next-state - the value of the state at time $(t+1)$ based on the present state and the input.
- Output - the value of the output as a function of the present state and (sometimes) the input.
- From the viewpoint of a truth table:
- the inputs are Input, Present State
- and the outputs are Output, Next State


## Example 1: State Table (from Fig. 5-15)

- The state table can be filled in using the next state and output equations:

$$
\begin{array}{ll}
- & \mathbf{A}(\mathbf{t}+\mathbf{1})=\mathbf{A}(\mathbf{t}) \mathbf{x}(\mathbf{t})+\mathbf{B}(\mathbf{t}) \mathbf{x}(\mathbf{t}) \\
& \mathbf{B}(\mathbf{t}+\mathbf{1})=\overline{\mathbf{A}}(\mathbf{t}) \mathbf{x}(\mathbf{t}) \\
- & \mathbf{y}(\mathbf{t})=\overline{\mathbf{x}}(\mathbf{t})(\mathbf{B}(\mathbf{t})+\mathbf{A}(\mathbf{t}))
\end{array}
$$



| Present State | Input | Next State |  | Output |
| :---: | :---: | :---: | :---: | :---: |
| $\mathbf{A}(\mathbf{t}) \mathbf{B}(\mathbf{t})$ | $\mathbf{x}(\mathbf{t})$ | $\mathbf{A}(\mathbf{t}+\mathbf{1})$ | $\mathbf{B}(\mathbf{t}+1)$ | $\mathbf{y}(\mathbf{t})$ |
| 0 | 0 | 0 | 0 | 0 |
| 0 |  |  |  |  |
| 0 | 0 | 1 | 0 | 1 |
| 0 | 1 | 0 | 0 | 0 |
| 0 | 1 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 | 0 |
| 1 | 0 | 1 | 1 | 0 |
| 1 | 1 | 0 | 0 | 0 |
| 1 | 1 | 1 | 1 | 0 |

## Example 1: Alternate State Table

- 2-dimensional table that matches well to a K-map. Present state rows and input columns in Gray code order.
- $\mathbf{A}(\mathbf{t}+\mathbf{1})=\mathbf{A}(\mathbf{t}) \mathbf{x}(\mathbf{t})+\mathbf{B}(\mathbf{t}) \mathbf{x}(\mathbf{t})$
- $\mathbf{B}(\mathbf{t}+\mathbf{1})=\overline{\mathbf{A}}(\mathbf{t}) \mathbf{x}(\mathbf{t})$
- $\mathbf{y}(\mathbf{t})=\overline{\mathbf{x}}(\mathbf{t})(\mathbf{B}(\mathbf{t})+\mathbf{A}(\mathbf{t}))$

| $\begin{gathered} \text { Present } \\ \text { State } \\ \text { A(t) B(t) } \\ \hline \end{gathered}$ | Next State |  | Output |  |
| :---: | :---: | :---: | :---: | :---: |
|  | $\mathrm{x}(\mathrm{t})=0$ | $\mathrm{x}(\mathrm{t})=1$ | $\mathbf{x}(\mathrm{t})=0$ | $\mathbf{x}(\mathrm{t})=1$ |
|  | $A(t+1) B(t+1)$ | $\mathbf{A}(\mathbf{t}+1) \mathbf{B}(\mathbf{t}+1)$ | $\mathrm{y}(\mathrm{t})$ | $\mathrm{y}(\mathrm{t})$ |
| 00 | 00 | 01 | 0 | 0 |
| 01 | 00 | 11 | 1 | 0 |
| 10 | 00 | 10 | 1 | 0 |
| 11 | 00 | 10 | 1 | 0 |

## State Diagrams

- The sequential circuit function can be represented in graphical form as a state diagram with the following components:
- A circle with the state name in it for each state
- A directed arc from the Present State to the Next State for each state transition
- A label on each directed arc with the Input values which causes the state transition, and
- A label:
- On each circle with the output value produced, or
- On each directed arc with the output value produced.


## State Diagrams

- Label form:
- On circle with output included:
- state/output
- Moore type output depends only on state
- On directed arc with the output included:
- input/output
- Mealy type output depends on state and input


## Example 1: State Diagram

- Which type?
- Diagram gets confusing for large circuits
- For small circuits, usually easier to understand than the state table



## Equivalent State Definitions

- Two states are equivalent if their response for each possible input sequence is an identical output sequence.
- Alternatively, two states are equivalent if their outputs produced for each input symbol is identical and their next states for each input symbol are the same or equivalent.


## Equivalent State Example

- Text Figure 5-17(a):
- For states S3 and S2,
- the output for input 0 is $\mathbf{1}$ and input $\mathbf{1}$ is 0 , and
- the next state for input 0 is $\mathbf{S 0}$ and for input 1 is $\mathbf{S 2}$.

- By the alternative definition, states S3 and S2 are equivalent.


## Equivalent State Example

- Replacing S3 and S2 by a single state gives state diagram:
- Examining the new diagram, states S1 and S2 are equivalent since
- their outputs for input 0 is 1 and input 1 is 0 , and
- their next state for input 0 is $\mathbf{S 0}$ and for input 1 is S 2 ,
- Replacing S1 and S2 by a single state gives state diagram:


## Moore and Mealy Models

- Sequential Circuits or Sequential Machines are also called Finite State Machines (FSMs). Two formal models exist:
- Moore Model
- Named after E.F. Moore
- Outputs are a function ONLY of states
- Usually specified on the states.
- Mealy Model
- Named after G. Mealy
- Outputs are a function of inputs AND states
- Usually specified on the state transition arcs.


## Moore and Mealy Example Diagrams

- Mealy Model State Diagram maps inputs and state to outputs

- Moore Model State Diagram maps states to outputs


## Moore and Mealy Example Tables

- Moore Model state table maps state to outputs

| Present | Next State |  | Output |
| :---: | :---: | :---: | :---: |
| State | $\mathbf{x}=0$ | $\mathbf{x}=1$ |  |
| 0 | 0 | 1 | 0 |
| 1 | 0 | 2 | 0 |
| 2 | 0 | 2 | 1 |

- Mealy Model state table maps inputs and state to outputs

| Present | Next State |  | Output |  |
| :---: | :---: | :---: | :---: | :---: |
| State | $\mathbf{x}=0$ | $x=1$ | $\mathbf{x}=0$ | $x=1$ |
| 0 | 0 | 1 | 0 | 0 |
| 1 | 0 | 1 | 0 | 1 |

## Mixed Moore and Mealy Outputs

- In real designs, some outputs may be Moore type and other outputs may be Mealy type.
- Example: Figure 5-17(a) can be modified to illustrate this
- State 00: Moore
- States 01, 10, and 11: Mealy
- Simplifies output specification



## Example 2: Sequential Circuit Analysis

- Logic Diagram:



## Example 2: Flip-Flop Input Equations

- Variables
- Inputs: None
- Outputs: Z
- State Variables: A, B, C
- Initialization: Reset to (0,0,0)
- Equations
- $\mathbf{A}(\mathbf{t}+\mathbf{1})=\mathbf{B C}$
- $\mathbf{B}(\mathbf{t}+\mathbf{1})=\mathbf{B}^{\prime} \mathbf{C}+\mathbf{B C}^{\prime}$
- $\mathbf{C}(\mathbf{t}+\mathbf{1})=\mathrm{A}^{\prime} \mathrm{C}^{\prime}$
- $\mathrm{Z}=\mathrm{A}$


## Example 2: State Table

## Example 2: State Diagram



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