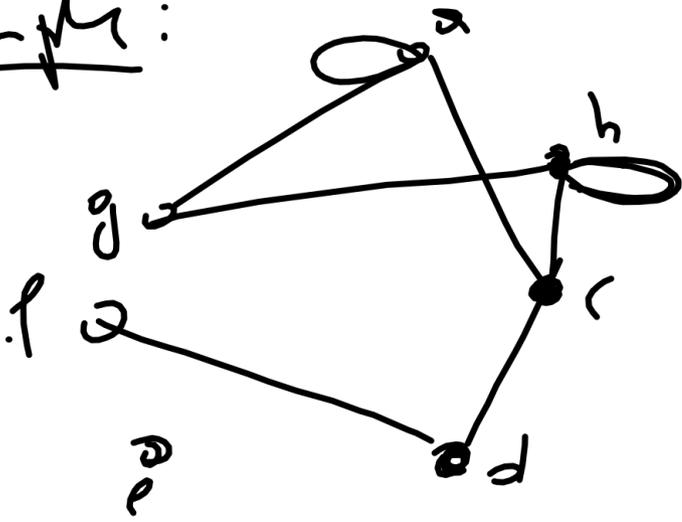


Example:



G

x	a	b	c	d	e	f	g
$d_G(x)$	4	4	3	2	0	1	2

$V = \{a, b, c, d, e, f, g\}$, $|V| = 7$
 e : isolated vertex ($d(e) = 0$)

f : Leaf or pendant vertex ($d(f) = 1$)

Sequence of degree of G or degree sequence of G
 $4, 4, 3, 2, 2, 1, 0 \rightarrow 7$ vertices.

$$\delta(G) = 0, \quad \Delta(G) = 4$$

$$d(G) = \frac{1}{n} \sum_{x \in V} \deg x = \frac{1}{7} (16) = \frac{16}{7}$$

average

we verify: $0 \leq \frac{16}{7} \leq 4$

$$\left(\delta(G) \leq d(G) \leq \Delta(G) \right)$$

Theorem: Let $G = (V, E)$ be a graph.

$$1) \sum_{x \in V} d(x) = 2|E|.$$

2) The number of odd vertices in V is even.

Example: $\sum_{x \in V} \deg(x) = 4 + 4 + 3 + 2 + 2 + 1 + 0$
 $= 16$

and we have $2|E| = 2(8) = 16$

2) Proof: Let $A = \{x: x \in V \text{ and } \deg(x) \text{ is even}\}$
 $B = \{x: x \in V \text{ and } \deg(x) \text{ is odd}\}$

we have $V = A \cup B$ and $A \cap B = \emptyset$

$$\sum_{x \in V} \deg(x) = \sum_{x \in A} \deg x + \sum_{x \in B} \deg x \stackrel{1)}{=} 2|E|$$

$$\text{Then } \sum_{x \in B} \deg x = 2|E| - \sum_{x \in A} \deg x$$

$\forall x: (x \in A \rightarrow \deg x \text{ is even})$

Then $\sum_{x \in A} \deg x$ is even and $2|E|$ is even
 $\underbrace{2|E|}_{\in \mathbb{N} \cup \{0\}}$

hence $2|E| - \sum_{x \in A} \deg x$ is even

Therefore $\sum_{x \in B} \deg x$ is even

Then $|B|$ is even
Then the number of odd vertices is even.

Proposition: Let $G = (V, E)$ be a graph.

$$\delta(G) \leq d(G) \leq \Delta(G)$$

$$\forall v \in V; \quad \delta(G) \leq \deg(v) \leq \Delta(G)$$

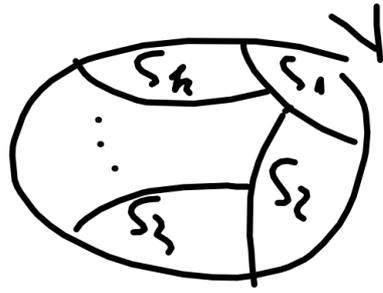
$$\text{Then } \sum_{v \in V} \delta(G) \leq \sum_{v \in V} \deg v \leq \sum_{v \in V} \Delta(G)$$

$$\text{hence } n \delta(G) \leq \sum_{v \in V} \deg v \leq n \Delta(G).$$

$$\text{Therefore } \delta(G) \leq \frac{1}{n} \sum_{v \in V} \deg v = d(G) \leq \Delta(G)$$

Pigeonhole Lemma

S_1, \dots, S_k is a partition of V , $|V| = n$
(i) $S_i \neq \emptyset, \forall i \in \{1, \dots, k\}$; (ii) $S_i \cap S_j = \emptyset, \forall i \neq j$; (iii) $V = \bigcup_{i=1}^k S_i$



Proposition: Let $G = (V, E)$ be a graph.

If G is a simple graph, then there are $x \neq y \in V$
such that $\deg x = \deg y$.

Proof: $V = \{x_1, x_2, \dots, x_n\}$; $|V| = n \geq 2$

By contradiction we suppose that, G is a simple graph
and $\forall x \neq y \in V$; $\deg x \neq \deg y$.

Suppose that $\deg x_1 < \deg x_2 < \dots < \deg x_n$

$$\sum_{x \in V} \deg x = \deg x_1 + \deg x_2 + \dots + \deg x_n = 2|E|$$

We have in G ; $\forall x \in V$; $0 \leq \deg x \leq n-1$

If $\deg x_1 = 0$, then $\deg x_n \leq n-2$ 

If $\deg x_n = n-1$, then $\deg x_1 \geq 1$

$\deg x \in \{0, 1, \dots, n-2\}$ or $\deg x \in \{1, \dots, n-1\}$.

where $|\{0, \dots, n-2\}| \leq n-1$ and $|\{1, \dots, n-1\}| \leq n-1$

$\exists x \neq y$; $\deg x = \deg y$

$\exists x \neq y$; $\deg x = \deg y$.

Exercises page 5:

1. (a) $S = (2, 3, 3, 4, 4, 5)$ sequence

Suppose that there is a graph $G = (V, E)$, such that

$S = (2, 3, 3, 4, 4, 5)$ is a sequence of degrees of G .

First method
The number of odd vertices (with $\deg v = 3$ or $\deg v = 5$)

is 3 (not even) contradiction

(with the number of odd vertices in graph is even)

Then there is no graph G with S as a sequence of degrees of G

Second method:

$$\sum_{v \in V} \deg v = 2 + 3 + 3 + 4 + 4 + 5 = 21 \neq 2|E| \text{ because } |E| \in \mathbb{N} \frac{1}{2}$$

(b) Suppose that there is a graph $G = (V, E)$
with $(2, 3, 4, 4, 4, 6, 6, 6, 9)$ as sequence of degrees.
We have $\forall x \in V; 0 \leq \deg x \leq n-1, n = |V|$.
In this case $n = |V| = 9$ and we have x_0 where
 $\deg x_0 = 9 > n-1 = 9-1 = 8$
is a contradiction.
Then there is no graph G , with S as the degree sequence of G .

(C) $S = (1, 3, 3, 3)$;

Suppose that there is a graph $G = (V, E)$

such that S is a degree sequence of G .

Let $a \in V$; where $\deg(a) = 1$, where $|V| = 4$

Suppose $V = \{a, v_1, v_2, v_3\}$; $\deg v_1 = \deg v_2 = \deg v_3 = 3$

Then $N(v_1) = \{a, v_2, v_3\}$; $N(v_2) = \{a, v_1, v_3\}$; $N(v_3) = \{a, v_1, v_2\}$

Then $\deg a = 3 > 1$ contradiction.

(d) Same case of (b)

(e) Same case of (a)

(f) $S = (2, 3, 4, 5, 5, 5)$

Suppose that there is a graph $G = (V, E)$ with S is sequence degree of G .

Let $x_1, x_2, x_3 \in V$; $\deg x_1 = \deg x_2 = \deg x_3 = 5$
 $|V| = 6$

Then $N(x_i) = V \setminus \{x_i\}$; $\forall i \in \{1, 2, 3\}$ ($\deg x_i = 5$)

Then $\delta(G) \geq 3$ is a contradiction with 2 is a degree of some vertex.

Then there is no a graph G , with S is a sequence degree of G .

2. Let $V = \{s_1, s_2, \dots, s_n\}$; V group of students
 $|V| = n \geq 2$

$x \neq y \in V$; ($xy \in E \Leftrightarrow x$ and y are friends).

$G = (V, E)$ is a graph.

Then there are $x \neq y \in V$; such that $\deg x = \deg y$

Then x and y are the same number of friends.

3. Let $V = \{c_1, c_2, \dots, c_{15}\}$

Let E be a set where

$x \neq y \in V; (xy \in E \Leftrightarrow x \text{ and } y \text{ are connected})$

Then $G = (V, E)$ is a graph.

Suppose that, is possible each computer connected with exactly 3 others.

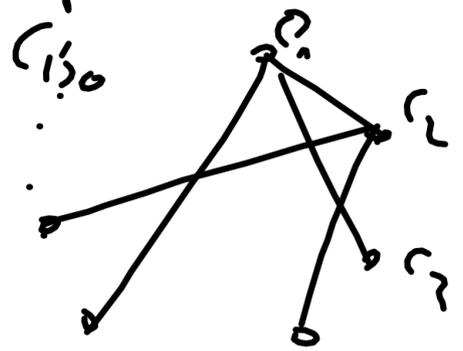
Then $\deg x = 3; \forall x \in V$

We have $\sum_{x \in V} \deg x = 2|E|$

hence $\sum_{x \in V} 3 = 2|E|$, then $(15) \cdot 3 = 2|E|$

then $|E| = \frac{45}{2} \notin \mathbb{N} \cup \{0\}$ (Contradiction)

Then there is not possible to connected each computer with exactly 3 others on 15 computers



4. The answer is same with question 3,
by replacing 3 by p and 1's by n ; $p < n$.

Exercise 9.1:

1) We have $\sum_{x \in V} \deg x = 2|E|$

Then $48 = 2|E|$, hence $|E| = 24$

2. Let $G = (V, E)$ be a regular graph with $\deg = 2$.

and $|V| = 14$

We have $\sum_{x \in V} \deg x = 2|E|$

Then $\sum_{x \in V} 2 = 2|E|$ because $\deg(x) = 2, \forall x \in V$

Since $|V| = 14$, then $2 \cdot 14 = 2|E|$, then for $|E| = 14$

3. Let $G = (V, E)$ be a graph with $|E| = 60$ and $|V| = n$.

$\sum_{x \in V} \deg x = 2|E|$, then $\sum_{x \in V} \deg x = 2(60) = 120$.

$\sum_{x \in V} \deg x = 120$

We compare with complete graph K_p

$(K_p = (W, F); |W| = p \text{ and } \forall x \neq y \in W, xy \in F)$

$|F| = \frac{p(p-1)}{2} \geq |E| = 60$

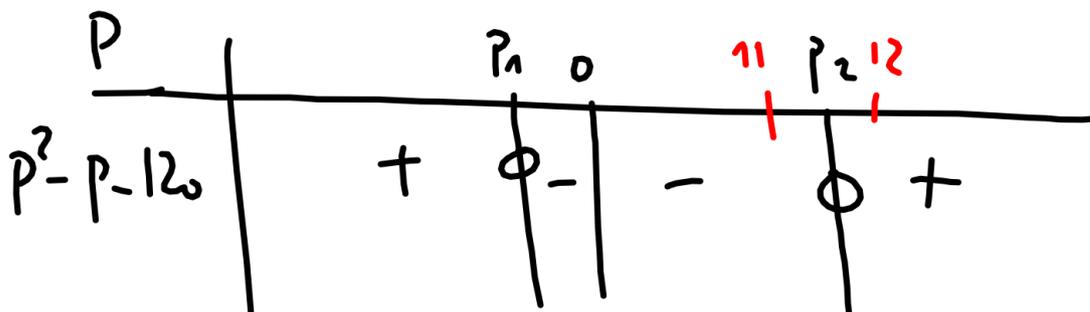
$(p^2 - p - 120 = 0)$

Then $p(p-1) \geq 120 \Leftrightarrow p^2 - p - 120 \geq 0$

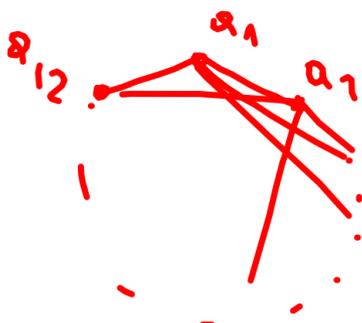
$\Delta = b^2 - 4ac = (-1)^2 - 4(1)(-120) = 1 + 480 = 481 =$

$p_1 = \frac{-b - \sqrt{\Delta}}{2a} = \frac{1 - \sqrt{481}}{2} < 0$

$p_2 = \frac{-b + \sqrt{\Delta}}{2} = \frac{1 + \sqrt{481}}{2} \approx 11.9..$



$p = 12 = n$



$G = K_{12} \{e_{11}, e_{12}, \dots, e_{12}, e_{11}\}$

Subgraph: Let $G = (V, E)$ and $H = (W, F)$ be two graphs.

- 1) H is a subgraph of G , if $W \subseteq V$ and $F \subseteq E$.
- 2) H is a spanning of G , if H is a subgraph of G and $W = V$.
- 3) Let $X \subseteq V$.

$H = (X, E \cap \{xy : x, y \in X\})$ is a subgraph of G
we said H is a induced subgraph of G on X , and
denoted by $H = \overline{G[X]}$

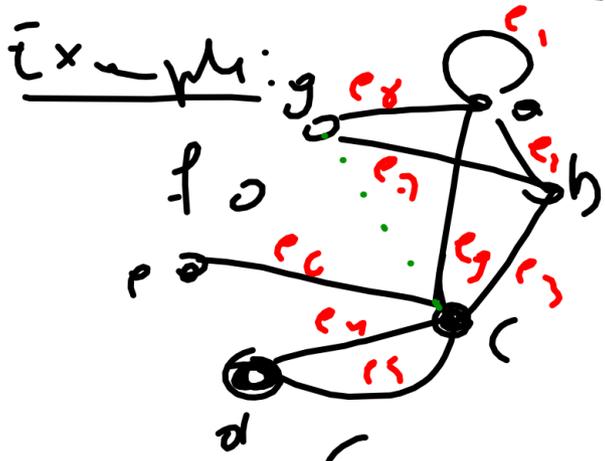
- 4) $H = G[V \setminus \{x\}]$; $x \in V$, is a subgraph induced
of G on $V \setminus \{x\}$, we denoted by $G - \{x\}$ or
 $G - x$

In general: Let $X \subseteq V$.

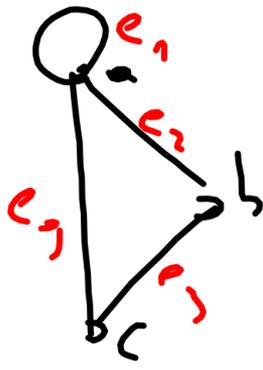
$H = G[V \setminus X]$ is a subgraph induced of G on
 $V \setminus X$, we denoted $H = G - X$

- 5) Let $E_1 \subseteq E$,
 $H = (V, E \setminus E_1)$ is a spanning graph of G .
we denoted $H = G - E_1$

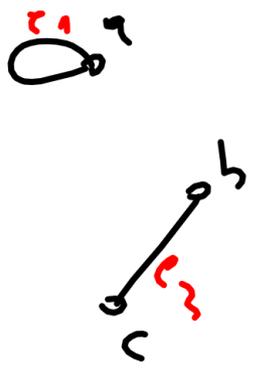
Subgraph: Let $G = (V, E)$ be a graph
 Let $H = (W, F)$ be a graph
 $H = (W, F)$ is a subgraph of G , if $W \subseteq V$ and $F \subseteq E$



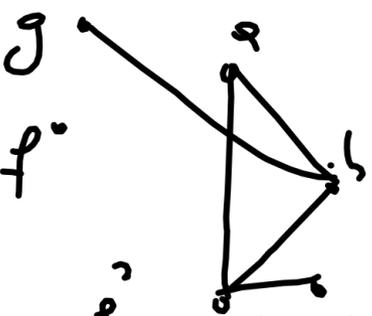
$V = \{a, b, c, d, e, f\}$
 $E = \{e_1, e_2, \dots, e_9\}$



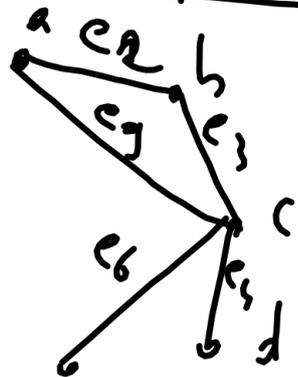
$H_1 = (W_1, F_1)$
 $W_1 = \{a, b, c\}$
 $F_1 = \{e_1, e_2, e_3, e_4\}$
 H_1 is a subgraph of G



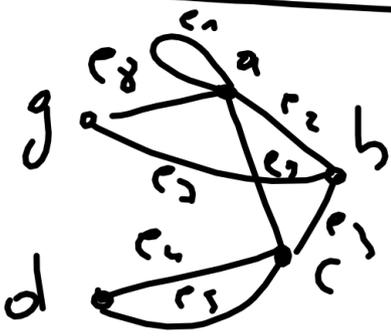
$H_2 = (W_2, F_2)$
 $W_2 = \{e, b, c, f\}$
 $F_2 = \{e_1, e_3\}$
 H_2 is a subgraph of G



$H_3 = (W_3, F_3)$ is not a subgraph of G
 because $a, d \notin E$ (then $F_3 \not\subseteq E$)



$H_4 = (V, E)$ is a subgraph of G
 H_4 is a spanning of G
 $H_4 = G - \{e_1, e_3, e_7, e_8\}$



$H_5 = G[\{a, b, c, d, g\}]$ induced subgraph of G on $\{a, b, c, d, g\}$
 $H_5 = G - \{e_1, f\}$

Walk-path: Let $G=(V, E)$ be a graph.

$$u, v \in V.$$

- A walk W of G from u to v , is a sequence x_0, x_1, \dots, x_p of V where $x_0 = u$ and $x_p = v$ and $x_i x_{i+1} \in E, i \in \{0, \dots, p-1\}$
- If $u=v$, A walk of G from u to v , we said w is a closed walk

- A trail of G from u to v , if a walk $w = (x_0 = u, x_1, \dots, x_p = v)$ where $\forall i \neq j; x_i x_{i+1} \neq x_j x_{j+1}, i, j \in \{0, \dots, p-1\}$.

- A closed trail or circuit $w = (x_0 = u, \dots, x_p = v)$ if w is trail and $u=v$

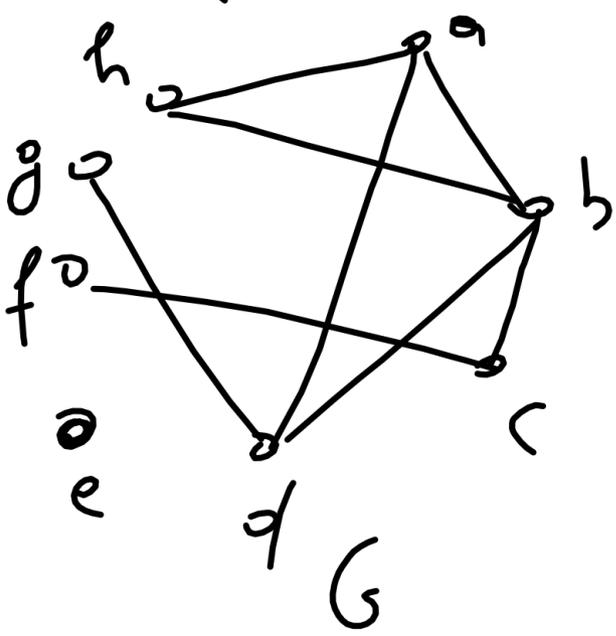
- A path P of G , from u to v , is a walk $P = (x_0 = u, x_1, \dots, x_p = v); x_i \neq x_j, \forall i \neq j$

- P is a closed path or cycle of G on u .

if $P = (x_0 = u, x_1, \dots, x_p = u)$ is a path of G from u to u .

- The length of w (walk of G), is the number of edges on w . denoted by $l(w) = p$

Example:



Walks of G from c to h.

- $W_1 = (c, b, a, h)$
- $W_2 = (c, f, c, b, d, a, h)$
- $W_3 = (c, b, a, d, b, h)$
- $W_4 = (c, h, a, d, b, h, a, h)$
- $W_5 = (c, b, h)$
- $W_6 = (c, h, a, b, h)$

Trail of G from c to h: $W_1, W_3, W_5,$

Path of G from c to h: W_1, W_5

Remark: Let G be a graph, and
let w be a walk

1) If w is a path, then w is a trail.

2) In general:

(If w is a trail, then w is a path) is false
for example in G : w is a trail, but not a path.

Proposition: Let G be a graph, let $x \neq y \in V(G)$.

If W is a walk from x to y in G , then

There is a path P contained in W , from x to y in G .

Proof: Let $W = x = x_0, x_1, \dots, x_p = y$ is a walk.

- If $x_i \neq x_j, \forall i \neq j$, then W is a path ($P = W$)
- If not, $\exists i \neq j$ such that $x_i = x_j$, W is not path.

→ Suppose that $i_0 \neq j_0, x_{i_0} = x_{j_0} \quad (i_0 < j_0)$

w.l.m. W is:
 $W = x_0 = x_1, x_2, \dots, x_{i_0} = x_{j_0}, \dots, x_{j_0+1}, \dots, x_p = y$.

Let $P_0: x_0 = x_1, x_2, \dots, x_{i_0} = x_{j_0}, x_{j_0+1}, \dots, x_p = y$



P_0 is a walk contained in W , from x to y

by iteration (p is finite), there is a path $P = x = x_0, y_1, \dots, y_k = y$

where $x_0, y_1, \dots, y_k \in \{x_0, x_1, \dots, x_p\}$

Then P is a path from x to y contained in W in G .

Definition: Let $G = (V, E)$ be a graph.
Let $x, y \in V$

1) x and y are connected, if there is a path from x to y or $x = y$, we denote $x \sim y$

2) G is connected graph, if $\forall x \neq y, x$ and y are connected

If not, G is disconnected graph.

Remark:

1) If x and y are adjacent, then x and y are connected

2) (if x and y are connected, then x and y are adjacent) is false, for example c and h in G (on the graph) are connected (with path $c \rightarrow t \rightarrow h$) and c and h are not adjacent.

Theorem: Let $G = (V, E)$ be a graph.

The relation on V defined by:

$$\forall x, y \in V; (x \mathcal{E} y) \equiv (x \text{ and } y \text{ are connected}).$$

We have \mathcal{E} is an equivalence relation on V .

Proof: Let $G = (V, E)$ be a graph.

- $x \mathcal{C} x$, because $x = x$, then \mathcal{C} is reflexive on V .
- Let $x, y \in V$, and suppose that $x \mathcal{C} y$, then there is a path \underline{P} from x to y in G .

When $\underline{P}: x = x_0, x_1, \dots, x_p = y$

then $\underline{P}^{-1}: y = x_p, x_{p-1}, \dots, x_1, x_0 = x$ is the inverse path of \underline{P}

and \underline{P}^{-1} is a path from y to x , then $y \mathcal{C} x$.

hence \mathcal{C} is a symmetric relation on V .

- Let $x, y, z \in V$; suppose that $x \mathcal{C} y$ and $y \mathcal{C} z$

then there is $\underline{P}_1: x = x_0, x_1, \dots, x_p = y$ a path from x to y

and there is $\underline{P}_2: y = y_0, y_1, \dots, y_k = z$ a path from y to z

hence $\underline{P} = \underline{P}_1 \cup \underline{P}_2: x = x_0, x_1, \dots, x_p = y = y_0, y_1, \dots, y_k = z$

is a walk from x to z by theorem,

there is a path \underline{Q} from x to z contained in \underline{P}

then $x \mathcal{C} z$

Therefore \mathcal{C} is transitive on V

Conclusion: \mathcal{C} is an equivalence relation on V .

Definition: Let $G = (V, E)$ be a graph.

• G is connected $\Leftrightarrow (\forall x, y \in V; \text{There is a path from } x \text{ to } y.)$

• \mathcal{C} : is an equivalence relation on V .

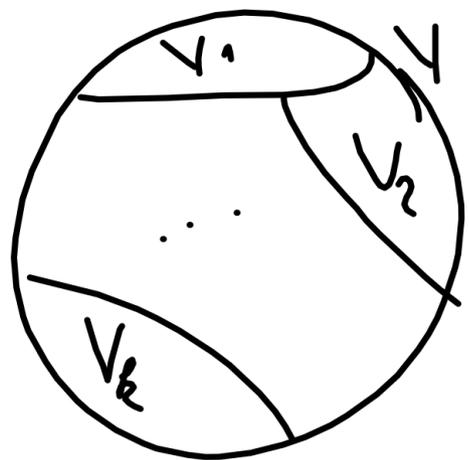
• In general case: Let $G = (V, E)$ be a graph.

There is a partition of V . (because \mathcal{C} is an equivalence relation)

where $V = V_1 \cup V_2 \cup \dots \cup V_p$, V_i is an equivalence class of \mathcal{C}

$\{V_1, V_2, \dots, V_p\}$ is a partition of V

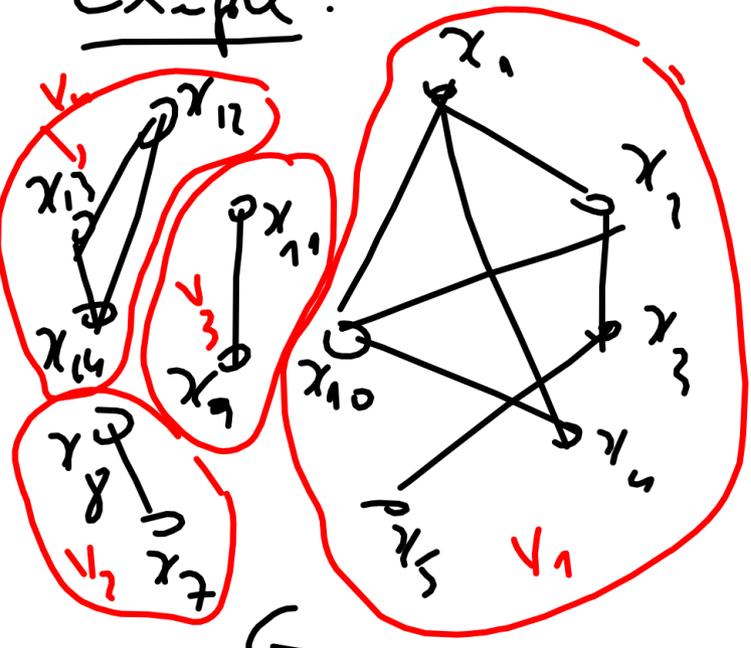
$\left(\begin{array}{l} \bullet \forall i \in \{1, \dots, p\}: V_i \neq \emptyset \\ \bullet \forall i \neq j: V_i \cap V_j = \emptyset \\ \bullet \bigcup_{i=1}^p V_i = V \end{array} \right)$



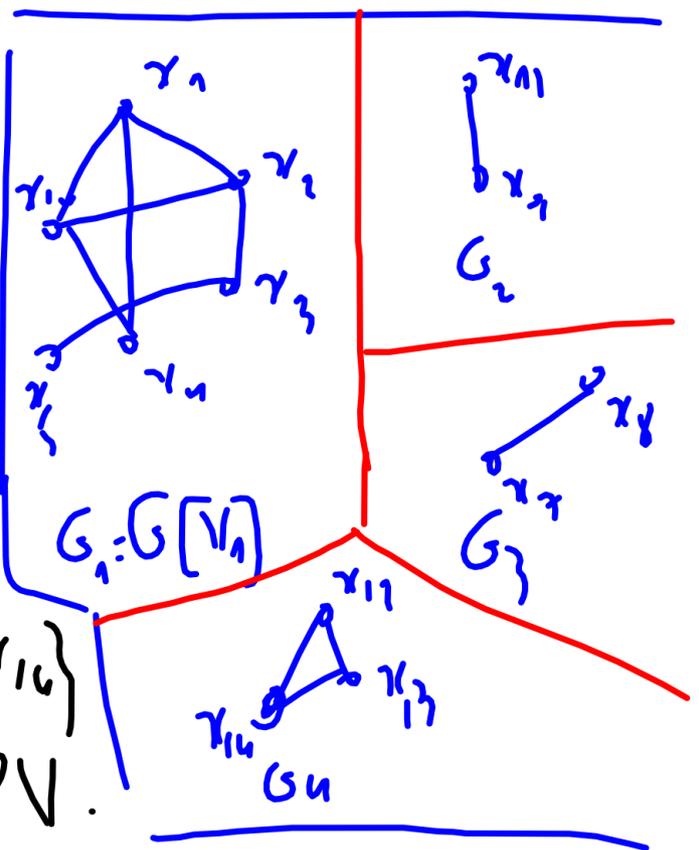
We said each V_i is a connected component of G
(and $G[V_i]; \forall i \in \{1, \dots, p\}$ is a connected component of G)

Remark: G is connected $(\Leftrightarrow) | G$ has only one connected component.
 $(\Leftrightarrow) (V \text{ is a connected component})$

Example:



$$V = \{x_1, x_2, \dots, x_{14}\}.$$



$$V_1 = \{x_1, x_2, x_3, x_4, x_5, x_{10}\}$$

$$V_2 = \{x_7, x_8\}$$

$$V_3 = \{x_9, x_{11}\}, V_4 = \{x_{12}, x_{13}, x_{14}\}$$

$\{V_1, V_2, V_3, V_4\}$ is a partition of V .

V_i is a connected component of G ; $\forall i \in \{1, 2, 3, 4\}$

$G[V_i]$ is a connected component of G ; $\forall i \in \{1, 2, 3, 4\}$

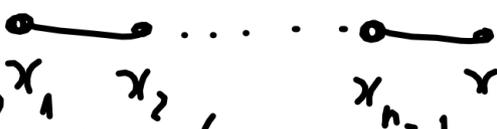
Particular Graph

1) Path and Cycles.

Let $G = (V, \bar{E})$ be a graph

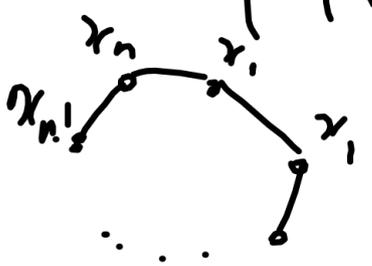
$$V = \{v_1, \dots, v_n\}; n \geq 1.$$

- G is a path if $V = \{x_1, x_2, \dots, x_n\}$ where
 $\bar{E} = \{x_i x_{i+1} : i \in \{1, \dots, n-1\}\}$



Proposition $x_1, x_2, \dots, x_{n-1}, x_n ; |V| = n$
 $DEG(G) = (1, 1, 2, \dots, 2), |E| = n-1$

- G is a cycle if $V = \{x_1, \dots, x_n\}$ where $n \geq 3$
 $E = \{x_i x_{i+1} : i \in \{1, \dots, n-1\}\} \cup \{x_1 x_n\}$, we send
 G is a cycle

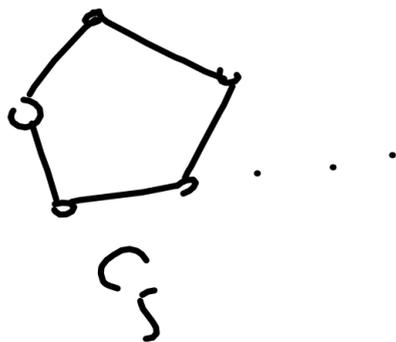
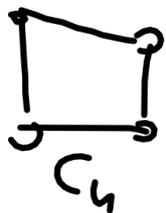
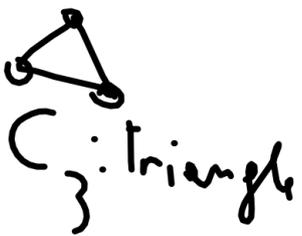
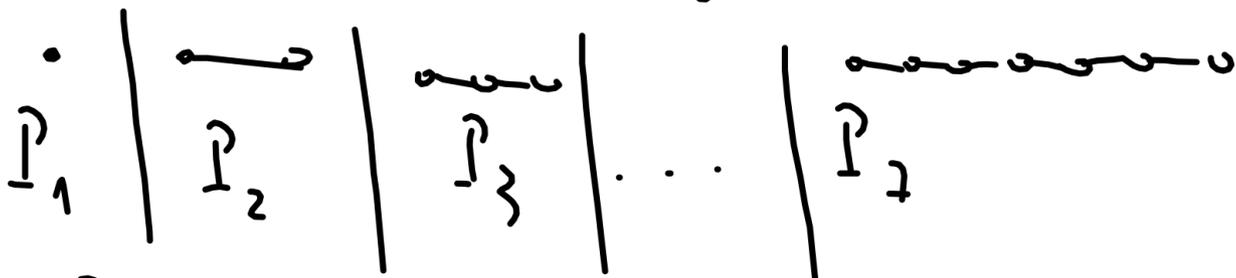


$|E| = |V| = n.$

$DEG(G) = (2, 2, \dots, 2).$

We denoted P_n : The path with n vertices

C_n : The cycle with n vertices



Let $G = (V, E)$ be a graph.

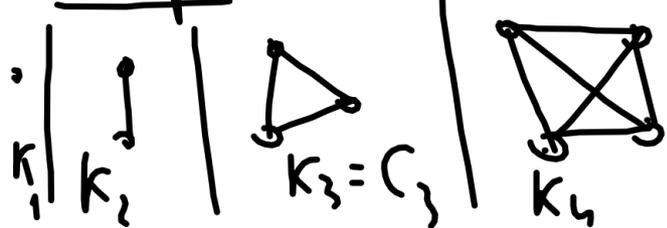
Complete graph:

G is a complete graph

if $\forall x \neq y \in V, xy \in E$

We denote $G = K_n$

Examp:



Null graph:

G is null graph on V with n vertices, if $E = \emptyset$
we denote D_n



Theorem: Let $G = (V, E)$ be a graph
If G is a complete graph, then
 $|E| = \frac{n(n-1)}{2}, |V| = n.$

$$|E| = 0.$$

Proof: Let $G = (V, E)$ be a graph

Suppose that G is a complete graph
with n vertices, then

$$E = \{xy : x \neq y \in V\}.$$

$$|E| = \binom{n}{2} = \frac{n!}{2!(n-2)!} = \frac{n \cdot (n-1)}{2}.$$

In K_n

$$\text{deg } x = n - 1$$

$$\forall x \in V$$

Regular graph: Let $G = (V, E)$ be a graph

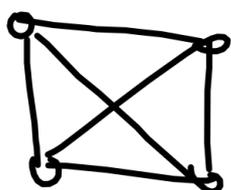
• G is a regular graph, if there is $r \in \mathbb{N} \cup \{0\}$ such that $\deg(v) = r$

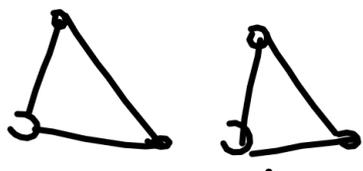
$$\text{DEG}(G) = (r, r, \dots, r)$$

In this case we said G is an r -regular graph.

Example:

1) The cycle is 2-regular graph

2) 
 G is 3-regular graph.

3) 
 G is 2-regular graph.

4) K_n is a $(n-1)$ -regular graph.

Theorem: Let $G = (V, E)$ be a graph, with $|V| = n$
If G is r -regular graph, then $|E| = \frac{n \cdot r}{2}$.

Proof: Assume that G is r -regular graph.

Then $\deg(x) = r, \forall x \in V$

We have: $\sum_{x \in V} \deg(x) = 2|E| \Rightarrow \sum_{x \in V} r = 2|E|$

hence $n \cdot r = 2|E| \Leftrightarrow |E| = \frac{n \cdot r}{2}$.

Bipartite graph:

Let $G = (V, E)$ be a graph.

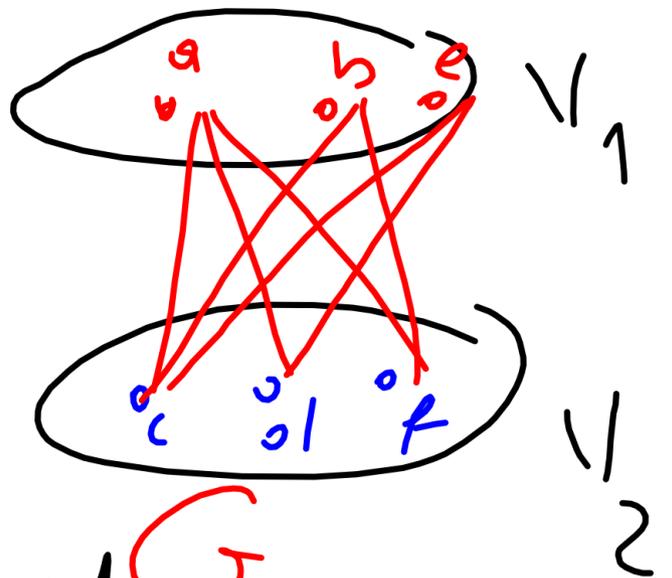
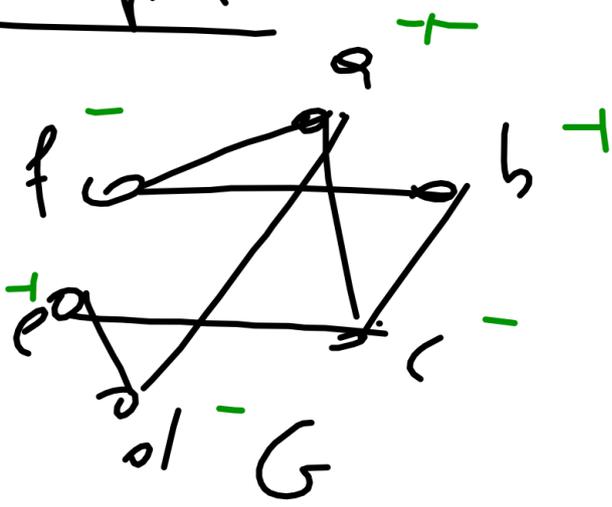
G is bipartite graph, if there is a bipartition $\{V_1, V_2\}$ on V .

($V = V_1 \cup V_2, V_i \neq \emptyset, \forall i \in \{1, 2\}$
 $V_1 \cap V_2 = \emptyset$)

such that $\forall x, y \in V_i$
 $xy \notin E$

We also note $G = (V_1 \cup V_2, E)$

Example:



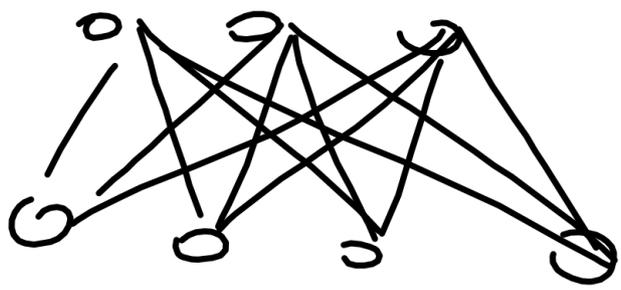
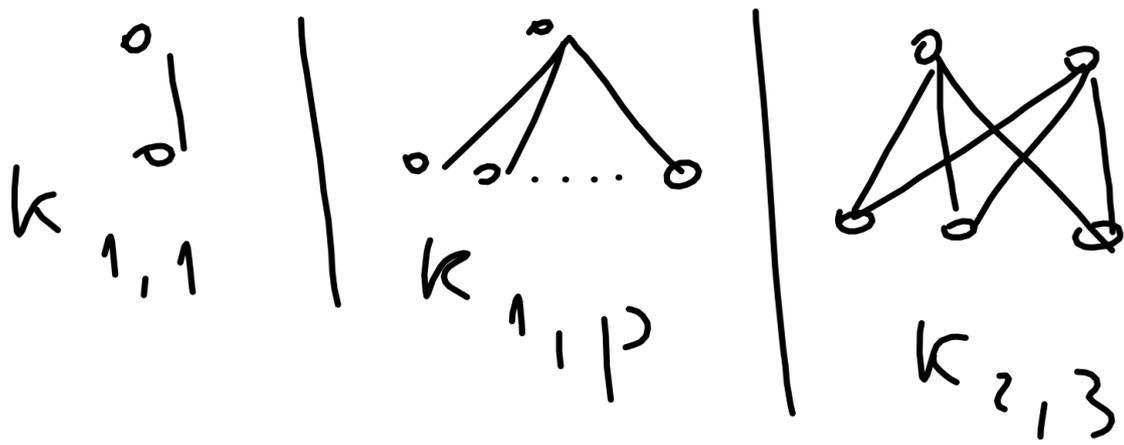
G is a bipartite graph. G

G is complete bipartite graph let $G = (V, E)$

if $G = (V_1 \cup V_2, E)$ is a bipartite graph and $\forall x \in V_1, \forall y \in V_2$
 $xy \in E$.

We denote $G = K_{n,m}$
where $n = |V_1|$ and $m = |V_2|$

Example:



$K_{3,4}$

Remark: $K_{1,p}$ is a star



Theorem: Let $G = (V, E)$ be a graph.

Let $n, m \in \mathbb{N}$

1) If $G = K_{n, m}$, then $|V| = n + m$

$$|E| = n \cdot m$$

2) $G = K_{n, m} = (V_1 \cup V_2, E)$; $|V_1| = n$

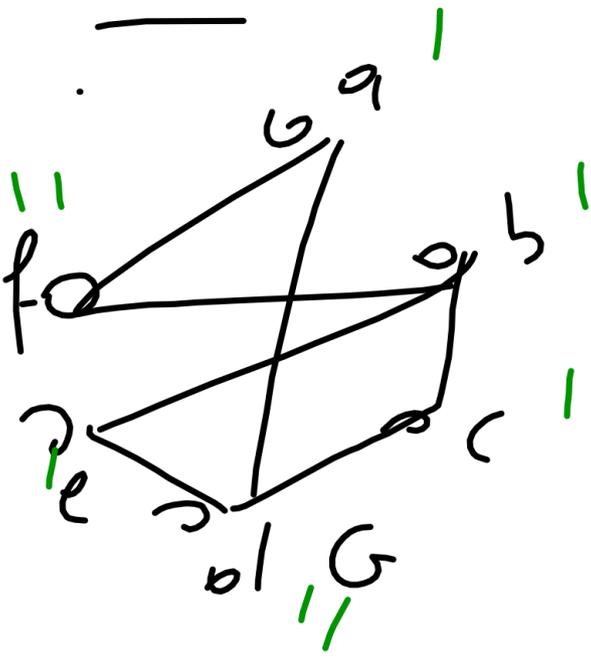
$\forall x \in V_1$, $\deg x = m$

$\forall x \in V_2$, $\deg x = n$.

3) G is bipartite graph if and only if G has not an odd cycle.

Remark: The odd cycle is not bipartite graph.

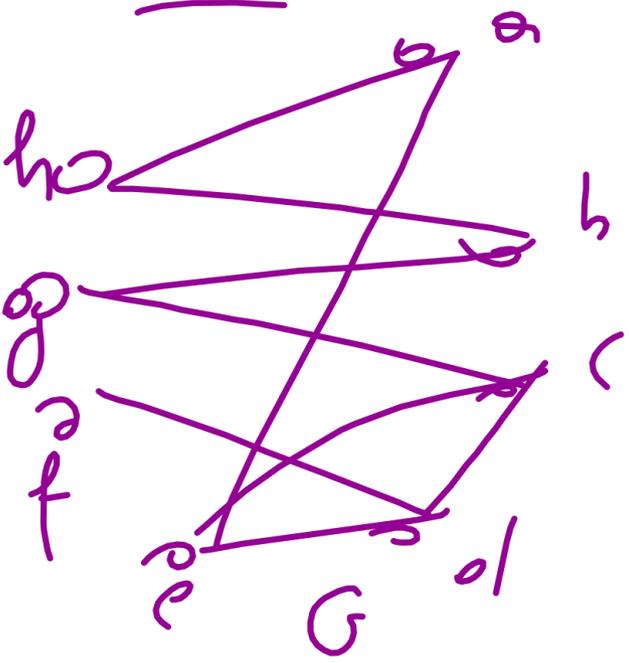
$G = \alpha$:



G is not bipartite graph, because

b, e, d, a, f, h is an odd cycle.
(5-cycle).

Ex:



G is not bipartite graph, because G has a 3-cycle (e, d, c, e)

Example: Is there a complete bipartite graph $G = (V_1 \cup V_2, E)$ such that $|V_1 \cup V_2| = 16$ and $|E| = 33$?

Solution: Suppose that there is a complete bipartite graph $G = (V_1 \cup V_2, E)$, where $|V_1| = n$ and $|V_2| = m$, we have:
 $n + m = 16$
 $n \cdot m = 33$.

• First method

$$n \in \{1, 3, 11, 33\}$$

$\rightarrow n = 1$, then $m = 16 - 1 = 15$ and $1 \cdot 15 \neq 33$

$\rightarrow n = 3$, then $m = 16 - 3 = 13$ and $3 \cdot 13 = 39 \neq 33$

$\rightarrow n = 11$, then $m = 16 - 11 = 5$, and $5 \cdot 11 = 55 \neq 33$

$\rightarrow n = 33$, then $m = 16 - 33 = \text{impossible}$

Conclusion:

then there is no a complete bipartite graph such that $|V| = |V_1 \cup V_2| = 16$ and $|E| = 33$.

• Second method:

$\{n, m\}$ is a solution of the equation

$$X^2 - SX + P = 0 \quad \text{where} \quad \begin{cases} n+m=S \\ n \cdot m=P \end{cases}$$

In this case: $X^2 - 16X + 33 = 0$

$$\Delta = b^2 - 4ac = (-16)^2 - 4 \cdot 1 \cdot 33 =$$

• Third method

n	1	2	3	4	5	6	7	8
$m = 16 - n$	15	14	13	12	11	10	9	8
$n \cdot m$	15	28	39	48	55	60	63	64

Isomorphism:

Def: Let $G = (V, E)$ and $H = (W, F)$ be two graphs.

G and H are isomorphic, if

there is ^{a function} f from V to W ($f: V \rightarrow W$)
a bijection such that

$$\forall x, y \in V \quad (xy \in E \iff (f(x)f(y)) \in F)$$

In this we said f is an isomorphism
from V to W (from G to H)

and we denoted $G \cong H$

$$\text{or } G \cong H \text{ or } G \cong H$$

$$\text{or } G \cong H$$

Theorem: Let $G = (V, E)$ and $H = (W, F)$

If G and H are isomorphic

Then 1) $|V| = |W|$ and $|E| = |F|$

2) $DEG(G) = DEG(H)$

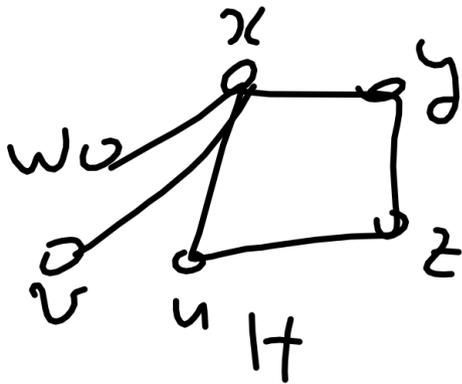
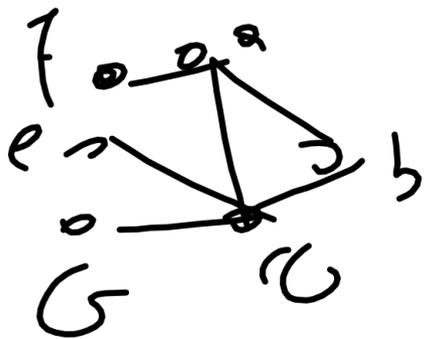
3) $\forall m \in \mathbb{N} \cup \{0\}$

$$\left| \left\{ x : \deg(x) = m \right\} \right|_{x \in V/G} = \left| \left\{ x : \deg(x) = m \right\} \right|_{x \in W/H}$$

Remark $(p \rightarrow q) \iff (q \rightarrow p)$

Remark: 1) $|E| \neq |F|$ or $|V| \neq |W|$, then $G \neq H$
2) If $\text{DEG}(G) \neq \text{DEG}(H)$
then $G \neq H$

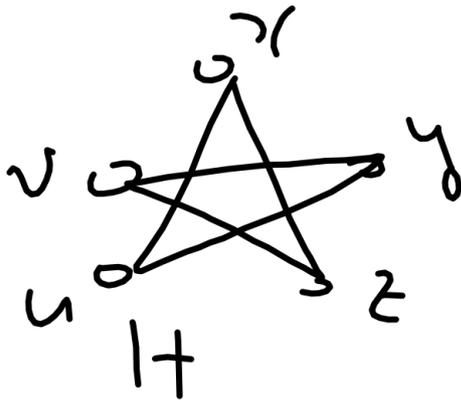
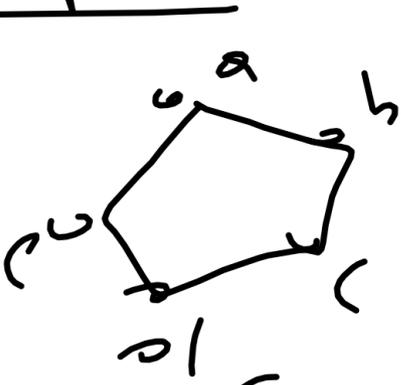
Example 1)



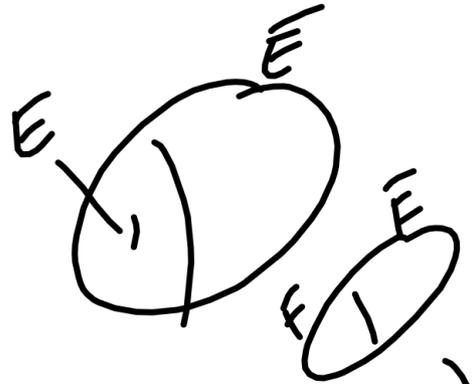
we have $|\{t : \deg t = 1\}| = 3$
 $\neq |\{t : \deg_H(t) = 1\}| = 2$

then G and H are not isomorphic.

Example 2:



Γ	a	b	c	d	e
$\Gamma(G)$	x	z	v	y	u



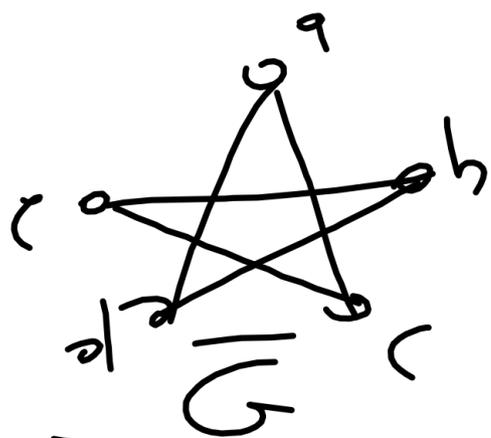
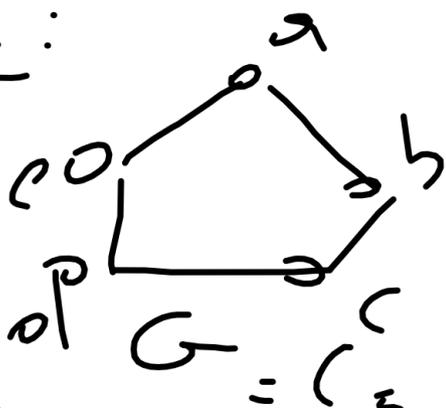
$|\Gamma(G)| = 5 = |\Gamma(H)|$
 $|\Gamma(G)| = \frac{3(5-1)}{2} - |\Gamma(G)| = 10 - 5 = 5$

- $a, b \in \Gamma(G) \rightarrow x, z \in \Gamma(H)$
- $b, c \in \Gamma(G) \rightarrow z, v \in \Gamma(H)$
- $c, d \in \Gamma(G) \rightarrow v, y \in \Gamma(H)$
- $d, e \in \Gamma(G) \rightarrow y, u \in \Gamma(H)$
- $e, f \in \Gamma(G) \rightarrow u, x \in \Gamma(H)$

Thus G and H are isomorphic
 (f is an isomorphism).

Def: Let $G = (V, E)$ be a graph.
 G is a self complementary graph,
if G and \bar{G} are isomorphic.
 $\bar{G} = (V, \bar{E})$: The complement graph of G .

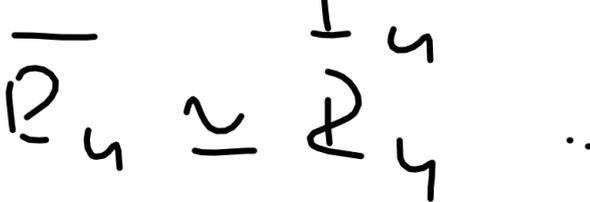
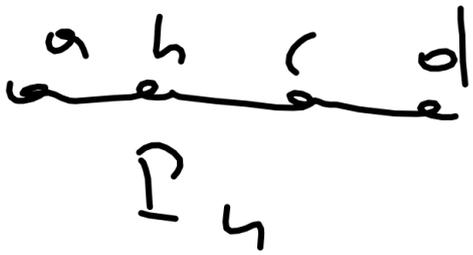
Example:



We know G and \bar{G} are isomorphic.
Then G is a self complementary
graph.

(C_3 is a self complementary graph.)

Example:



Then P_4 is a self complementary graph.

Theorem: Let $G = (V, E)$ be a graph

We have G or \bar{G} is connected

(Rh: It is not possible to find a graph

Proof: G and \bar{G} are both disconnected)

• If G is connected, then we have a result.

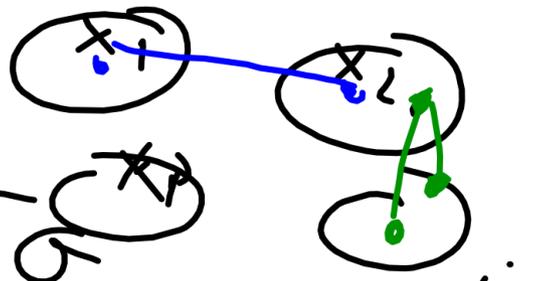
• If G is disconnected then they are p connected components



$X_1, \dots, X_p; p \geq 2$

$\rightarrow x \in X_i, y \in X_j, i \neq j$

Since $xy \in \bar{E}(G)$
Then $xy \in E(\bar{G})$



\rightarrow Let $i \in \{1, \dots, p\}; x \neq y \in X_i$

$\forall j \neq i \in \{1, \dots, p\}$, let $z \in X_j$
 $xz \in \bar{E}(G)$ and $zy \in \bar{E}(G)$

hence x, z, y is a path on \bar{G}

hence \bar{G} is connected graph.

conclusion G or \bar{G} is connected graph.

Theorem: Let $G = (V, E)$ be a graph with $|V| = n$

If G is a self-complementary graph, then

1) G and \bar{G} are connected graphs.

2) $|V| = \begin{cases} 4p & p \in \mathbb{N} \\ 4p+1, & p \in \mathbb{N} \cup \{0\} \end{cases}$

Proof: We know G or \bar{G} are connected (for a reason then not).

Without loss of generality, we suppose that G is connected, since G and \bar{G}

are isomorphic (G is self-complementary) then \bar{G} is connected.

Therefore G and \bar{G} are connected.

2) Let G is self-complementary graph then $|\bar{E}(G)| = |\bar{E}(\bar{G})|$

and we have $|\bar{E}(G)| + |\bar{E}(\bar{G})| = \frac{n(n-1)}{2}$

then $2|\bar{E}(G)| = \frac{n(n-1)}{2}$

hence $|\bar{E}(G)| = \frac{n(n-1)}{4} \in \mathbb{N} \cup \{0\}$

\times $\frac{n}{4} = 4p, \frac{n(n-1)}{4} \in \mathbb{N} \cup \{0\}$

$n = 4p+1$, then $n-1 = 4p$, hence $\frac{n(n-1)}{4} \in \mathbb{N}$

$n = 4p+2$, then $\frac{(4p+2)(4p+1)}{4} = \frac{(2p+1)(4p+1)}{2} \notin \mathbb{N}$

$n = 4p+3$, then $\frac{n(n-1)}{4} = \frac{(4p+3)(4p+2)}{4} \notin \mathbb{N}$

$= \frac{(4p+1+2)(2p+1)}{2} \notin \mathbb{N} \cup \{0\}$

hence $n = \begin{cases} 0 \\ 1 \\ 4p \\ 4p+1 \end{cases}$ (mod 4)

Exercise: Prove that:

1) C_n is a self complementary graph if and only if $n=5$.

2) P_n is a self complementary if and only if

Proof: 1) " \Rightarrow " C_n is self complementary.

$$n=4p \text{ or } n=4p+1; p \in \mathbb{N} \cup \{0\}$$

$$|E(C_n)| = |E(\overline{C_n})|$$

$$\downarrow \quad \downarrow$$

$$n \quad \cdot \quad \frac{n(n-1)}{2} - n = \frac{n^2 - 2n}{2}$$

$$n^2 - 2n = 2n \Leftrightarrow n^2 - 4n = 0$$

$$\Leftrightarrow n(n-4) = 0$$

$$\Leftrightarrow n=0 \text{ or } n=4 \text{ (} n \neq 0 \text{)}$$

" " \Leftarrow " $\Rightarrow n=4$

Let $n=4$ we have (in example)

$C_4 \cong \overline{C_4}$, hence C_4 is self complementary

2) " \Rightarrow " P_n is self complementary,

$$\text{then } |E(P_n)| = |E(\overline{P_n})|$$

$$\text{then } n-1 = \frac{n(n-1)}{2} - (n-1)$$

$$\Leftrightarrow (n-1)(n-4) = 0$$

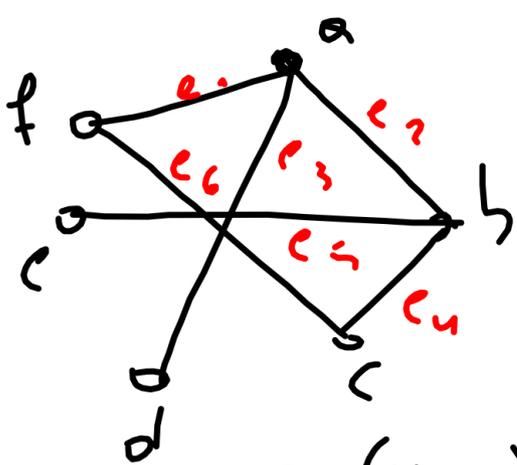
$$\Leftrightarrow n=1 \text{ or } n=4$$

" " \Leftarrow " P_4 is self complementary (in example)

P_1 is a graph.

then P_1 (resp. P_4) is self complementary

Adjacency matrix and length of path



Adjacency matrix of G

$$A(G) = \begin{matrix} & \begin{matrix} a & b & c & d & e & f \end{matrix} \\ \begin{matrix} a \\ b \\ c \\ d \\ e \\ f \end{matrix} & \begin{bmatrix} 0 & 1 & 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 \end{bmatrix} \end{matrix}$$

$$G = (V, E)$$

$$V = \{a, b, c, d, e, f\}$$

Incidence matrix

$$I(G) = \begin{matrix} & \begin{matrix} e_1 & e_2 & e_3 & e_4 & e_5 & e_6 \end{matrix} \\ \begin{matrix} a \\ b \\ c \\ d \\ e \\ f \end{matrix} & \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 \end{bmatrix} \end{matrix}$$

Theorem: Let $G = (V, E)$ be a graph.

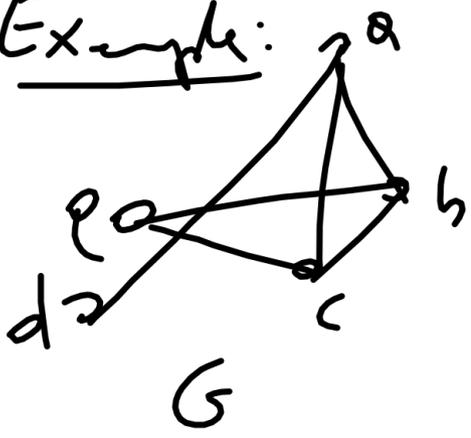
$V = \{v_1, \dots, v_n\}$; $A = (a_{ij})$ the adjacency matrix of G ;
$$a_{ij} = \begin{cases} 1 & \text{if } v_i v_j \in E \text{ ; } i \neq j \\ 0 & \text{if } v_i v_j \notin E \end{cases} \quad 1 \leq i, j \leq n$$

Let $k \in \mathbb{N}$

$$A^k = (b_{ij})_{1 \leq i, j \leq n}$$

b_{ij} : The number of walks from v_i to v_j with length k

Example:



$$A = \begin{matrix} & \begin{matrix} a & b & c & d & e \end{matrix} \\ \begin{matrix} a \\ b \\ c \\ d \\ e \end{matrix} & \begin{bmatrix} 0 & 1 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 & 1 \\ 1 & 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 & 0 \end{bmatrix} \end{matrix}$$

How many walks from a to d, with length 4.

$$M^2 = \begin{bmatrix} 0 & 1 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 & 1 \\ 1 & 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 & 0 \end{bmatrix} \cdot \begin{bmatrix} 0 & 1 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 & 1 \\ 1 & 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 & 0 \end{bmatrix} = \begin{bmatrix} 3 & 1 & 1 & 0 & 2 \\ \dots & \dots & \dots & \dots & \dots \\ \dots & \dots & \dots & \dots & \dots \\ \dots & \dots & \dots & \dots & \dots \\ \dots & \dots & \dots & \dots & \dots \end{bmatrix}$$

$$M^4 = M^2 \cdot M^2 = \begin{bmatrix} 3 & 1 & 1 & 0 & 2 \\ \dots & \dots & \dots & \dots & \dots \\ \dots & \dots & \dots & \dots & \dots \\ \dots & \dots & \dots & \dots & \dots \\ \dots & \dots & \dots & \dots & \dots \end{bmatrix} \cdot \begin{bmatrix} 3 & 1 & 1 & 0 & 2 \\ \dots & \dots & \dots & \dots & \dots \\ \dots & \dots & \dots & \dots & \dots \\ \dots & \dots & \dots & \dots & \dots \\ \dots & \dots & \dots & \dots & \dots \end{bmatrix} = \begin{bmatrix} \dots & \dots & \dots & \dots & \dots \\ \dots & \dots & \dots & \dots & \dots \\ \dots & \dots & \dots & \dots & \dots \\ \dots & \dots & \dots & \dots & \dots \\ \dots & \dots & \dots & \dots & \dots \end{bmatrix}$$

There is 2 walks from a to d with length 4.

- a b c a d → not path | trail
- a c b a d → not path | trail

There is 0 paths from a to d with length 4.

There is 2 trails from a to d